## Outline

Greedy MST

## Minimum spanning tree



## Minimum Spanning Tree

Given $G=(V, E), w: E \rightarrow \mathbb{R}$, a minimum spanning tree $T$ is a spanning tree (i.e. connecting all vertices) that minimizes $\operatorname{cost}(T)=\sum_{e \in T} w(e)$


## Cut Property

## Lemma

Let $T$ be a minimum spanning tree, $X \subset T$ s.t. $X$ does not connect $(S, V-S)$. Let $e$ be the lightest edge from $S$ to $V-S . X \cup e$ is part of some MST.


## Cut Property Proof

## Cut Property

Let $T$ be an MST, $X \subset T$ s.t. $X$ does not connect $(S, V-S)$. Let $e$ be the lightest edge from $S$ to $V-S . X \cup e$ is part of some MST.

Let $X \subseteq T$ where $T$ is MST
$>$ if $e \in T$, done
$>$ add $e$ to $T$, makes a cycle

## Cut Property Proof

- Let $X \subseteq T$ where $T$ is MST
- if $e \in T$, done
- add $e$ to $T$, makes a cycle

- $\exists$ crossing $e^{\prime} \in E-X$
- swap $e$ and $e^{\prime}$


## Prim's Algorithm


$S=$ nodes reached so far

## Prim's Algorithm

def prim(G,root):
$\mathrm{pq}=\mathrm{pqdict}() ; \operatorname{prev}=\{ \}$
for $v$ in G.keys ():
pq.additem(v,inf)
pq.updateitem (root, 0)
while len(pq) >0:
$\mathrm{v}=\mathrm{pq} \cdot \mathrm{pop}()$
for (z,weight) in G[v]:
if $z$ in $p q$ and weight < pq[z]:
prev[z]=v
pq.updateitem(z,weight)
return prev

```
0 \inA
0 \in V-A
```



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0 \inA
0 \in V-A
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0 \inA
0 \in V-A
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```
o \inA
0 \in V-A
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## ALGORITHMS <br> Example of Prim's algorithm

```
0 \inA
\bullet \in V-A
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## ALGORITHMS <br> Example of Prim's algorithm

```
0 }\in
- \(\in V-A\)
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o \inA
- \(\in V-A\)
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\(0 \in A\)
- \(\in V-A\)
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\(0 \in A\)
- \(\in V-A\)
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- \(\in A\)
- \(\in V-A\)
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\(0 \in A\)
- \(\in V-A\)
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- \(\in A\)
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