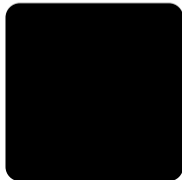


# CS3383 Unit 0, Lecture 1: Deeper into Asymptotics

David Bremner David Bremner



## Asymptotics

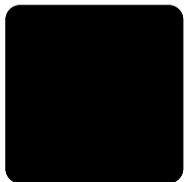
big-Theta = big-O *and* big-Omega

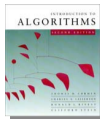
Root vs. Log

An example we didn't get to

Introducing little-o and little-omega

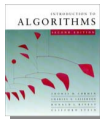
Coming back to our example





## $\Theta$ -notation (tight bounds)

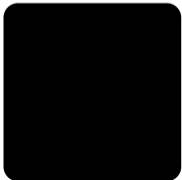
$$\Theta(g(n)) = O(g(n)) \cap \Omega(g(n))$$



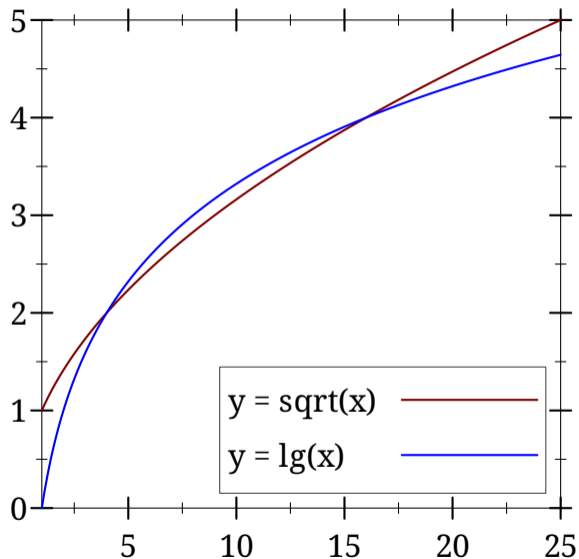
## $\Theta$ -notation (tight bounds)

$$\Theta(g(n)) = O(g(n)) \cap \Omega(g(n))$$

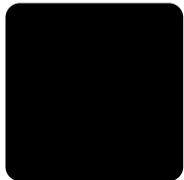
**EXAMPLE:**  $\frac{1}{2}n^2 - 2n = \Theta(n^2)$



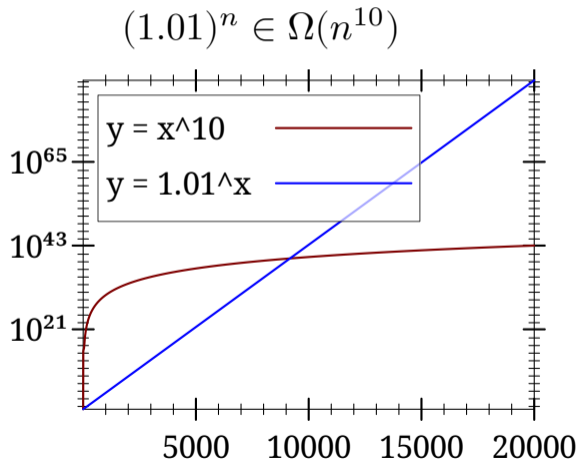
## Root vs lg, revisited



► note there are 2 crossing points



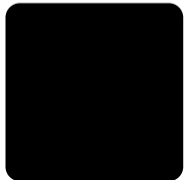
# Exponential versus Polynomial



## CLRS3.13

$$\lim_{n \rightarrow \infty} \frac{n^b}{a^n} = 0 \quad \forall a > 1$$

- ▶ How to prove?
- ▶ How does it help?



# Strengthening big-O

Definition ( $f(n) \in O(g(n))$ )

$$\exists c \exists n_0 \quad \forall n > n_0 \quad 0 \leq f(n) \leq cg(n)$$

Definition ( $f(n) \in o(g(n))$ )

$$\forall c > 0 \exists n_0 \quad \forall n > n_0 \quad 0 \leq f(n) < cg(n)$$

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# Strengthening big-O

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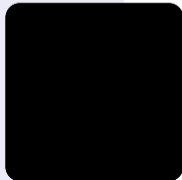
example

$f \in o(g) \implies f \in O(g)$ , but not vice-versa.

$$2n^2 \in o(n^3)$$

$$2n^2 \in O(n^2)$$

$$2n^2 \notin o(n^2)$$



# Strengthening big-Omega

Definition ( $f(n) \in \Omega(g(n))$ )

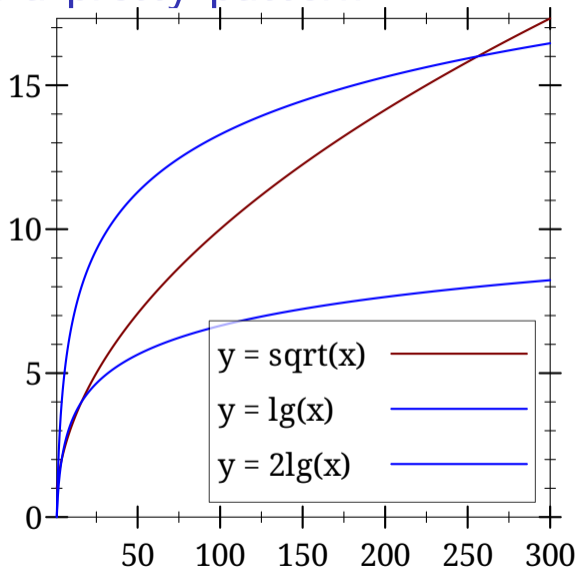
$$\exists c \exists n_0 \quad \forall n > n_0 \quad 0 \leq cg(n) \leq f(n)$$

Definition ( $f(n) \in \omega(g(n))$ )

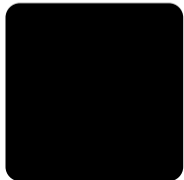
$$\forall c > 0 \exists n_0 \quad \forall n > n_0 \quad 0 \leq cg(n) < f(n)$$



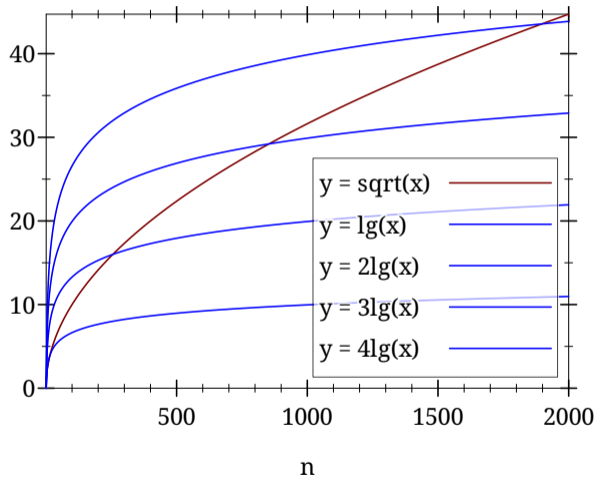
## Ooh a pretty pattern



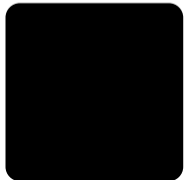
- ▶ Want:  
 $\sqrt{n} \in \omega(\lg(n))$
- ▶ For  $c = 1$ ,  $n_0 = 16$   
(last class)
- ▶ For  $c = 2$ ,  $n_0 = 256$



# The pattern is a lie



- ▶ For  $c = 3$ ,  
 $n_0 \approx 853.25$
- ▶ For  $c = 4$ ,  
 $n_0 \approx 1897.41$ .
- ▶ The exact solution, is  
not *nice*



# Sometimes calculus is the easy way...

Definition ( $f(n) \in \omega(g(n))$ )

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \infty$$

A useful rule (CLRS3.13)

For any  $b$ , and any  $a > 1$

$$\lim_{n \rightarrow \infty} \frac{a^n}{n^b} = \infty$$

